

What is *i*KP?

- *i*-Key-Protocol, *i* = 1, 2, 3, ...
- Family of protocols
- Secure electronic payment
- Based on credit card payments
- Can be extended to debit card and check payments

History of *i*KP

- 1995, IBM Research Labs Zurich and Watson Research Centre
- Open industry standard
- Incorporated into SEPP, SET
- ZiP: fully operational prototype
- Did not become commercial product
- But has been deployed in some businesses

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Initial Assumptions

- Only partial privacy is emphasized
- Encryption not used in protocol
- Can be implemented by other means
- Or protocol can be extended
- *i*KP emphasizes the payment
- Assumes purchase order is already known

Parties and Attackers

- Three parties: Acquirer (A), Merchant (M), Customer (C)
- Three attackers: eavesdropper, active attacker, insider

Acquirer Requirements

- A1. Proof of transaction authorization of customer
- A2. Proof of transaction authorization by merchant

Merchant Requirements

- M1. Proof of transaction authorization by acquirer
- M2. Proof of transaction authorization by customer

Customer Requirements

- C1. Unauthorized payment is impossible
- C2. Proof of transaction authorization by acquirer
- C3. Certification and authentication of merchant
- C4. Receipt from merchant

Extra Customer Requirements C5. Privacy C6. Anonymity

Components of the Protocol

- Keys: PK_x, SK_x, CERT_x
- Cryptography: H(.), E_x(.), S_x(.)
- Quantities: SALT_c, PRICE, DATE, NONCE_M, ID_M, TID_M, DESC, CAN, R_c, CID, Y/N, PIN, V
- Discuss 1KP in detail, comparison with 2KP and 3KP

1KP Protocol Flow

Initiate: $C \rightarrow M$ Invoice: $C \leftarrow M$ Payment: $C \rightarrow M$ Auth-Request: $M \rightarrow A$ Auth-Response: $M \leftarrow A$ Confirm: $C \leftarrow M$

SALT_c, CID Clear EncSlip Clear, H(DESC,SALT_c), EncSlip Y/N, Sig_A Y/N, Sig_A

Removing Nonce and Date

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Initiate: C→M
Invoice: C←M
Payment: C→M
Auth-Request: M→A
Auth-Response: M←A
Confirm: C←M
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SALT_C, CID Clear EncSlip* Clear, H(DESC,SALT_C), EncSlip* Y/N, Sig_A Y/N, Sig_A

 $\begin{array}{l} \mbox{Common: } \mbox{PRICE, } \mbox{ID}_{M}, \mbox{TID}_{M}, \mbox{DATE, NONCE}_{*}, \mbox{Clear: } \mbox{ID}_{M}, \mbox{TID}_{M}, \mbox{DATE, NONCE}_{*}, \mbox{H(Common)} \\ \mbox{SLIP: } \mbox{PRICE, } \mbox{H(Common), } \mbox{CAN, } \mbox{R}_{C}, \mbox{[PIN]} \\ \mbox{EncSlip: } \mbox{E}_{A}(SLIP) \\ \mbox{Sig}_{A}, \mbox{S}_{A}(Y/N, \mbox{H(Common)}) \\ \mbox{Sig2004} \\ \mbox{Sig204} \\ \mbox{Sig204$

Removing SALT_C

Initiate: C→M Invoice: C←M Payment: C→M Auth-Request: M→A Auth-Response: M←A Confirm: C←M

Clear EncSlip Clear, H(DESC,SALT_c), EncSlip Y/N, Sig_A Y/N, Sig₄

 $\begin{array}{l} \mbox{Common: } \mbox{PRICE, } ID_M, \mbox{TID}_M, \mbox{DATE, } NONCE_M, \mbox{CID, } H(\mbox{DESC, SALT}_c) \\ \mbox{Clear: } ID_M, \mbox{TID}_M, \mbox{DATE, } NONCE_M, \mbox{H}(\mbox{Common}) \\ \mbox{SLIP: } \mbox{PRICE, } H(\mbox{Common}), \mbox{CAN, } R_c, \mbox{[PIN]} \\ \mbox{Clear: } \mbox{Clear: } D_M, \mbox{Cl$ EncSlip: $E_A(SLIP)$ Sig_A: $S_A(Y/N, H(Common))$

Removing SALT_C: Invariant

-- Intruder never knows desc and salt from same customer invariant "Intruder does not know DESC" forall i: IntruderId do forall j: CustomerId do !int[i].descs[j] !int[i].salts[j] end end; 3/12/2004

What is not fulfilled?

- Acquirer's proof of transaction authorization by merchant
- Merchant's proof of transaction authorization by customer
- Customer's certification and authentication of merchant
- Customer's receipt from merchant

2KP and 3KP

- Satisfies deficiencies of 1KP
- Differs in the number of public keys available
- Guarantees more undeniable receipts for the transaction

2KP Protocol Flow

Initiate: C→M Invoice: C←M Payment: C→M Auth-Response: M←A Y/N, Sig_A Confirm: C←M

SALT_C, CID Clear, $\underline{Sig}_{\underline{M}}$, $\underline{CERT}_{\underline{M}}$ EncSlip Auth-Request: $M \rightarrow A$ Clear, H(DESC,SALT_c), EncSlip, <u>Sig</u>_M, <u>CERT</u>_M

Y/N, <u>V</u>, Sig_A

 $\begin{array}{l} \mbox{Common: PRICE, ID_M, TID_M, DATE, NONCE_M, CID, H(DESC, SALT_C), \underline{H(V)} \\ \mbox{Clear: ID_M, TID_M, DATE, NONCE_M, \underline{H(V)}, H(Common) \\ \mbox{SLIP: PRICE, H(Common), CAN, R_C} \\ \mbox{EncStip: E_A(SLIP) \\ \mbox{EncStip: E_A$ Sig_A: S_A(Y/N, H(Common)) Sig_M: S_M(H(Common), H(V)) 3/12/2004

2KP Additions

- Merchant has public/private key and certificate
- Acquirer has certification and authentication of merchant
- Customer has certification and authentication of merchant
- Customer has receipt from merchant

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3KP Protocol Flow

Initiate: C→M Invoice: C←M Payment: C→M

SALT_C, CID, CERT_C Clear, <u>Sig</u>_M EncSlip, Sigc

 Auth-Request: $M \rightarrow A$ Clear, $H(DESC,SALT_c)$, EncSlip, $\underline{Sig_M}$, $\underline{Sig_c}$

 Auth-Response: $M \leftarrow A$ Y/N, Sig_A

 Confirm: $C \leftarrow M$ Y/N, \underline{V} , Sig_A

Common: PRICE, ID_M, TID_M, DATE, NONCE_M, CID, H(DESC,SALT_c), <u>H(V)</u> Clear: ID_M, TID_M, DATE, NONCE_M, <u>H(V)</u>, H(Common) SLIP: PRICE, H(Common), CAN, R_c EncSlip: E_A(SLIP) Sig_A: S_A(YN, H(Common)) Sig_A: S_A(H(Common), <u>H(V)</u>) Sig_A: S_A(H(Common), <u>H(V)</u>)

3KP Additions

- Customer has public/private key and certificate
- Merchant has proof of transaction authorization by customer

iKP Comparison Overview

Requirements / Protocols	1KP	2KP	3KP
A1. Proof of Transaction Authorization by Customer	*	*	**
A2. Proof of Transaction Authorization by Merchant		**	**
M1. Proof of Transaction Authorization by Acquirer	**	**	**
M2. Proof of Transaction Authorization by Customer			**
C1. Unauthorized Payment is Impossible	*	*	**
C2. Proof of transaction Authorization by Acquirer	*	*	**
C3. Certification and Authentication of Merchant		**	**
C4. Receipt from Merchant		**	**

*i*KP Summary

- 1KP: simple, merchant is not authenticated, order and receipt are deniable
- 2KP: merchant is authenticated
- 3KP: non-repudiation for all messages
- Intended for gradual deployment

Bibliography

- Bellare, Mihir *et al.* Design, Implementation and Deployment of the *i*KP Secure Electronic Payment System. IEEE Journal of Selected Areas in Communications, VOL 18, NO. 4, April 2000.
- Bellare, Mihir *et al. i*KP A Family of Secure Electronic Payment Protocols. 1995.